#### **Outline: exam review**

- List, ArrayList, LinkedList
- parametrized types
- iterators
- invariants
- reminders: recursion, runtime analysis, binary search

#### exam review

- format similar to last exam
- material from lecture notes (including in-class exercises), book, assignments, quizzes
- for the book, all the material in Chapters 2-2.9 and 5.
- must also be familiar with the material presented before exam 1, i.e. review the material (and the exam, quizzes, and homework assignments)
  - classes and interfaces, runtime analysis, binary search.
- review all the code posted on the course web page. Understand this code well enough to be able to code similar programs
  - the actual question may ask for something similar, and the corresponding code may be the same or different.

## **Collection and List**

- hierarchy of parametrized interfaces and classes
  - including abstract classes and non-abstract (concrete) classes
  - List and Collection are interfaces, ArrayList and LinkedList are classes
- the type parameter E indicates the types of objects stored by the collection

interface List<E> extends Collection<E> { ...

• elements of a list are in order and have an index

### List

- variable sized collection of objects in a particular sequence
  - <sup>–</sup> being in sequence, each element corresponds to a specific index in 0..list.size()-1
- duplicate and null elements are permitted
  - unless the list is special-purpose, e.g. in some of the assignments
- the List interface defines required methods
- the AbstractList class helps in implementing new list classes
- some particularly useful methods of List:
  - boolean add(E e) add at the end of a list
  - void add(int index, E e) add within the list
  - Iterator<E> iterator() create a new iterator
  - E remove(int index) remove and return the object at the given position
  - int size()

# ArrayList

- a list in which the elements are stored in an array, usually called data
- the array may have more elements than the list, so an instance variable (size) is needed to keep track of the size of the list
  - capacity() returns data.length, the size of the underlying array
  - invariant: size <= capacity()</pre>
- as the list grows, it may need a new array, so in the worst case adding at the end of the array is O(n)
  - as long as the array grows by doubling the capacity, the time for adding n elements is O(n), meaning adding at the end takes <u>amortized</u> constant time
- adding and removing at the beginning or in the middle of the array is always O(n) -make sure you understand why

# LinkedList

- a list in which each element is stored in a node, and nodes are linked to each other
- the number of nodes is exactly the same as the number of elements
- each node has a reference to the value stored and a reference to the next node in the list: item or value, and next
- the end of the list is reached when the next field has the value null
- a linked list must have as instance variable the head of the list
- adding at the end of the linked list is faster if a tail pointer is kept -- O(1) instead of O(n)
- adding and removing at the front of a linked list is always O(1)

# Circular and doubly-linked list

- in a circular list, the next field of the last LinkedNode refers (back) to the head of the list
- in a doubly-linked list, each node has the next reference and also a reference to the node before it (prev or previous)
  - the previous node of the head is null, and the next node of the tail is null
- in a doubly-linked circular list, the tail is the previous node for the head, and the head is the next node for the tail
  - prev and next are never null

#### Iterators

- an iterator is an object that keeps track of the state of a traversal of a collection
- an iterator is like a bookmark for a book: there can be many bookmarks for a given book
  - more than one of them might be on the same page
- calling the iterator() method of a collection class is the common way to create a new iterator
- once the iterator exists, calls to hasNext() and next() may continue until there are no more elements
  - or, can repeat calling next() until it throws an exception
- examples:

#### **Using Iterators: Example 1/3**

- calls to hasNext() and next() may continue until there are no more elements
- for example:

```
List<E> list = ...
Iterator<E> iter = list.iterator();
while (iter.hasNext()) {
  E variable = iter.next();
    // can use "variable" in the loop
```

#### **Using Iterators: Example 2/3**

• can repeat calling next() until it throws an exception

```
List<E> list = ...
Iterator<E> iter = list.iterator();
while (true) {
 try {
    E variable = iter.next();
      // can use "variable" in the loop
  } catch (NoSuchElementException e) {
    break;
  }
```

## Using Iterators: foreach (Example 3/3)

 as an alternative to explicitly calling the iterator methods, use the foreach style:

```
for (E variable: list) {
    // can use "variable" in the loop
}
```

list must implement Iterable<E> (or be an array)

### **Building an Iterator**

- for the exam, must understand and be able to implement at least the hasNext() and next() methods of LinkedListIterator.java and an iterator for ArrayLists, and of other simple iterators
- each iterator is its own separate class
  - and may be internal to a collection class
- the iterator needs access to the elements of the collection
  - perhaps provided to the constructor
  - perhaps from the instance variables of the enclosing collection class
- each iterator needs its own variables to keep track of its <u>state</u>, that is, where it is in the iteration

### Invariants

- an invariant is something the programmer believes to be true about relationships among the program's variables
- for example: in binary search, if the object we are looking for is in the array, it <u>must</u> be between indices start..end
  - this requires that the array must be sorted using the same comparison operation that is used in the binary search
- in an object-oriented language such as Java, the most important relationships are among an object's instance variables (class data fields)
- if a public method is called with the invariant being true, the invariant <u>must</u> still be true at the end of the call
- a class invariant <u>must</u> always be true at the end of a public constructor

#### recursion

- recursion is useful when we have a problem that:
  - has an easy solution for some base cases, and
  - for all other cases, has a solution that can be expressed in terms of solving a problem that is <u>closer to the base case</u>.
- the problem that is <u>closer to the base case</u> is often a smaller problem
  - e.g., smaller value of n
  - e.g., linked list shorter by one node
- recursion is a way of thinking as well as a programming technique
  - mathematicians often think recursively

#### recursion examples on linked lists

 the length of a linked list beginning with node is one more than the length of the list beginning with node.next:

```
private long listLength(LinkedNode<E> node) {
    if (node == null) return 0;
    return 1 + listLength(node.next);
}
```

• adding to or removing from a linked list is easy if we are willing to create new LinkedNodes for the backbone of the linked list:

```
private LinkedNode<E> remove(LinkedNode<E> node, int index) {
    if (node == null) return null;
    if (index == 0) return node.next; // remove this value
    return new LinkedNode<E>(node.value, remove(node.next, index - 1));
```

called as follows: head = remove(head, index);

## runtime analysis

- constant time: linked list insertion or deletion at the front, linked list insertion at the back if a tail pointer is kept
  - <u>amortized</u> constant time for insertion at the end of an array list
- log time if doubling the size of the problem increases the time by a constant: binary search
  - later we will see tree operations that take log time
- linear time: ordered list insertion or deletion, linear search
- quadratic time: removing n elements from the end of a singly-linked list or the front of an array list
- exponential time if adding one element doubles the time: towers of Hanoi, clear but inefficient computation of fibonacci numbers