Outline

- equality and comparisons in Java
- sorting
- selection sort
- bubble sort
- insertion sort
- introduction to algorithm analysis

Java Equality details for Objects

- four kinds of equality:
- == is true if and only if two references are to the exact same object
 - for basic types, if the values are the same, e.g. x == 3 is true if x is an int with value 3
- object1.equals(object2) is true if object1's equals method decides (in its wisdom) that they are the same
- object1.compareTo(object2) is 0 if the compareTo method decides (in its wisdom) that they are the same
 - the result is negative if object1 < object2, and greater than 0 if object1 > object2)
- object0.compare(object1, object2) is similar to compareTo, comparing object1 to object2
- the behavior of the last three depends on the implementation:
 - ideally, they compare the contents of the object
 - ideally they are all consistent with each other
 - but:
 - they may not do what you think
 - they may have bugs

properties of the equals method

the equals method should be:

- reflexive: a.equals(a) should always be true
- symmetric: a.equals(b) == b.equals(a)
 - both true or both false
- transitive: (a.equals(b) && b.equals(c)) only if a.equals(c)
- consistent: successive identical calls should return the same result
- if a is not null, a.equals (null) should always be false
- the equals method can be overridden for any class
- if equals is not overridden, Object.equals is the same as ==

ordering objects in Java

- two Java methods to find out which object is greater and which object is less:
- Comparable<T> interface, specifies the int compareTo(T x) method
- Comparator<T> interface, specifies the int compare (T x1, T x2) method (which could be static, but isn't interfaces don't list static methods)
- if comparing a to b:
 - that is, if we are calling x.compareTo(y) or compare(x, y)
 - both methods return:
 - a value that is < 0 if x < y
 - 0 if x equals y
 - a value that is > 0 if x > y
 - these methods are not part of the Object class, so are not provided by all objects

compare and compareTo

- given that signum is the sign function, returning -1 for negative numbers, 0 for zero, and 1 for positive numbers:
- both comparison methods should have the following properties
 - symmetric:

Integer.signum(compare(a, b)) == - Integer.signum(compare(b, a))

- transitive: if (((compare(a, b) > 0) && ((compare(b, c) > 0)),
 - then ((compare(a, c) > 0))
- consistent: if (compare(a, b) == 0), then the sign of (compare(a, c)) should be the same as the sign of (compare(b, c))
- when sorting, can use either interface
- e.g. class Collections has two different methods for sorting:
 - static <T extends Comparable<? super T>> void sort(List<T> list) works when elements are Comparable
 - static <T> void sort(List<T> list, Comparator<? super T> c) takes as parameter a Comparator
- compare and compareTo are essentially the same, but some classes only provide one of the two

Applications of sorting

- finding duplicate elements in a list or an array
 - or eliminating duplicates by making a list of elements, where each element occurs at most once
- preparing to make future searches more efficient
 - dictionary, phone book
- presenting data in an appropriate format, e.g. for printing
 - transactions in a bank statement are sorted by date
- comparing two lists to find out which elements are in one, the other, or both
- merging multiple sorted collections into a new sorted collection, with or without duplicates

Sorting with the Java Standard Library

- Java class Arrays has
 public static void sort(x[] items);
 public static void
 sort(x[] items, int fromIndex, int toIndex);
- for type <u>x</u> being any one of:
- int and all of the basic types except boolean
- Object the objects must be Comparable
- T if the last parameter is a Comparator<T>

Selection Sort

- start with an unsorted array a
- find the smallest element in the array, at index *i*
- swap the element at index *i* with the element at index 0
- now, find the smallest element in 1..a.length-1
- swap the elements at index I' and at index 1
- now, the sub-array with elements 0..1 is sorted
- now, find the smallest element in 2..a.length 1
 - see animations at https://en.wikipedia.org/wiki/Selection_sort
- in-class exercise (in groups of 2 or 3): write code to implement selection sort

Bubble Sort

- start with an unsorted array a
- loop through all the elements of array a, swapping any that are not sorted
- repeat until the array is sorted
- see animations at https://en.wikipedia.org/wiki/Bubble_sort
- trivia: bubble sort of an array of size n can be done in parallel on n processors

Insertion Sort

- in selection sort, find the smallest element, and put it in the next position
- in <u>insertion sort</u>, take the next element, and put it in the right place
- trivia: card players often use insertion sort to arrange their decks
- the sub-array at the beginning of the array is already sorted, just as in selection sort
- in selection sort, the next element e is always taken from the next index in the array
- elements in the sorted part of the array that are greater than e are shifted up one position, to free the position where element e will be stored
 - see animations at https://en.wikipedia.org/wiki/Insertion_sort
- in-class exercise (in groups of 2 or 3): write code to implement insertion sort

Introduction to Algorithm Analysis: motivation

- selection sort, bubble sort, and insertion sort all sort an array
- which one of these is better?
 - assuming the implementations already exist
- all of these algorithms sort <u>in place</u>, so they all use the same amount of memory space
- so the only difference may be in the time they take
- but the time depends on both the computing system and the data
 - for example, bubble sort is very fast on data that is already sorted, or if sorting data stored on magnetic tape

Introduction to Algorithm Analysis: analysis

- for each algorithm, we will have a <u>best-case time</u>, a <u>worst-case</u> <u>time</u>, and sometimes an <u>average-case time</u>
- the time taken for sorting depends on the size of the array: sorting 1000 elements takes longer than sorting 10 elements
- because computing systems vary, and to keep things simple, all constant-time operations are assumed to take the same amount of time
 - so this analysis assumes that addition, multiplication, and comparisons, each take the same, constant time

Introduction to Algorithm Analysis: notation

- If, even in the worst case, an algorithm takes constant time, we say it is O(1)
 - for example, the time to access an element of an array is constant, no matter the size of the array
- If the worst case of an algorithm takes time proportional to the size of the input, it is O(n), with <u>n</u> the size of the input
 - in the worst case, finding out whether an unsorted array of integers contains a specific number (called <u>linear</u> <u>search</u>) takes time proportional to the size of the array
 - linear search in an array of 1,000,000 elements should be twice as fast as in an array of 2,000,000 elements
- If an algorithm in the worst case takes time proportional to the square of n, we say it is O(n²)
- We will also see O(log n), O(n log n), O(\sqrt{n}), and so on
- O() always looks at the <u>worst case</u> as n gets very large
 - unless we are specifically looking for the best case or average case