

Outline

- reminder: invariants
- **Big-O Analysis of array list efficiency**
- **Linked Lists**

Reminder: Invariants

- a useful property of a program that the programmer knows and that is always true, is an *invariant*
 - in the ArrayList implementation one invariant is that the elements of data from 0..size-1 are all of type E
 - the invariant must be true whenever a public method is called and when a public method returns, but may not be true in the middle of a method body
 - every invariant must hold by the end of each constructor, that is, constructors must **establish** all the invariants
 - for example, the ArrayList class has the invariant that $\text{size} \leq \text{data.length}$
 - in the ArrayList `add`, we may change size before we assign the new value: the invariant is temporarily broken when size is changed, then restored before the method completes

Reminder: Lists

- a list holds a sequence of values
- each value is at a specific index in the list
- we can insert items into a list and remove items from a list
 - insert or remove at the beginning, at the end, or at a specific index

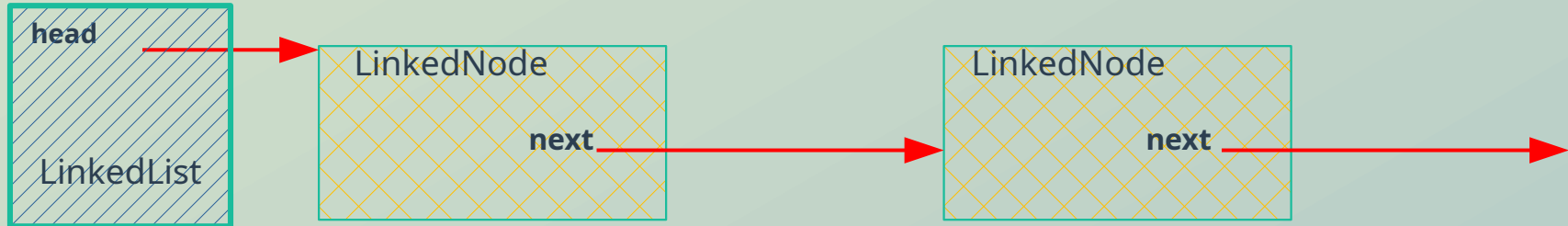
Linked Lists

and comparison with Array Lists

- in an array list, the elements of the list are stored in the array
- if more elements are needed, the array is replaced by a bigger array
- that means the elements are stored in a single object (an array) whose size may be different depending on the number of elements
- instead, consider storing the elements into a variable number of fixed-size objects
- where each object can only store one element
- there will be as many such objects as elements in the list

Linked List Nodes

- objects used to store list elements are called Nodes
- a **Linked List** has zero or more nodes
 - so there is a `LinkedListNode` class and a `LinkedList` class
- how to keep track of all these nodes?
- each of the nodes has a reference (*link*) to another node



- the `LinkedList` class has one variable to keep track of the first element, the *head* of the linked list

Linked Lists Nodes in Java

```
private class ListNode<E> {  
    private E item;                // value in this node  
    private ListNode<E> next;      // reference to the next node, or null  
  
    private ListNode(E value) {    // constructor  
        item = value;  
        next = null;              // last node in the linked list  
    }  
  
    private ListNode(E value, ListNode<E> reference) { // constructor  
        item = value;  
        next = reference;  
    }  
}
```

Linked Lists Details

- each node has two data fields: an element value, and a reference to the next node
- the reference to the next node is a pointer to the next node:
 - the `next` reference of the last node in the linked list must be `null`
 - every other `next` reference in the linked list must refer to the next node

Linked Lists Implementation: class structure

- **LinkedListNode is a private class inside LinkedList**
- **a private class is local to the enclosing (parent) class**
- **all data fields (meaning, all instance variables) in the private class are accessible to the code in the parent class**
 - even private instance variables are accessible to the code in the parent class!

LinkedList Implementation

- the linked list class only really needs one class variable: a reference to the start of the list, conventionally known as the head of the list
- the book also has a size class variable, which can be useful for error checking
- my implementation (and most other implementations) also keeps a reference to the last node in the linked list, the tail of the list

In-Class Exercises

- for a linked list:
- what is the run-time performance of `add(E)`?
- what is the run-time performance of `add(int, E)`?
- what is the run-time performance of `remove()`, which removes the head of the linked list?
- what is the run-time performance of `remove(int)`, which removes an element at an arbitrary position?