Overview of Storage & Indexing (i)

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Data Storage

• Main Memory
  – Random access
  – Volatile
• Flash Memory
  – Random access
  – Random writes are expensive
• Disk
  – Random access
  – Sequential access cheaper
• Tapes
  – Only sequential access
  – Archiving
Relational Tables on Disk

- **Record** -- a tuple or row of a relational table
- **RIDs** – record identifiers that uniquely identify a record across memory and disk
- **Page** – a collection of records that is the unit of transfer between memory and disk
- **Bufferpool** – a piece of memory used to cache data and index pages.
- **Buffer Manager** – a component of a DBMS that manages the pages in memory
- **Disk Space Manager** – a component of a DBMS that manages pages on disk
Magnetic Disks

- A disk or platter contains multiple concentric rings called **tracks**.
- Tracks of a fixed diameter of a spindle of disks form a **cylinder**.
- Each track is divided into fixed sized **sectors** (i.e. “arcs”).
- Data stored in units of disk **blocks** (in multiples of sectors)
- An array of **disk heads** moves as a single unit.
- **Seek time**: time to move disk heads over the required track
- **Rotational delay**: time for desired sector to rotate under the disk head.
- **Transfer time**: time to actually read/write the data
Accessing Data on Disk

- **Seek time**: time to move disk heads over the required track
- **Rotational delay**: time for desired sector to rotate under the disk head.
  - Assume uniform distribution, on average time for half a rotation
- **Transfer time**: time to actually read/write the data
Example: Barracuda 1TB HDD (ST31000528AS)

• What is the average time to read 2048 bytes of data?
  
  = Seek time + rotational latency + transfer time

  = $8.5 \text{ msec} + 4.16 \text{ msec} + \left( \frac{2048}{512} \right) / 63 \times \left( \frac{60000 \text{ msec}}{7200 \text{ rpm}} \right)$
  
  = $8.5 + 4.16 + 0.265$

<table>
<thead>
<tr>
<th>Feature</th>
<th>Value</th>
</tr>
</thead>
<tbody>
<tr>
<td>cylinders</td>
<td>121601</td>
</tr>
<tr>
<td>Bytes/cylinder</td>
<td>16065*512</td>
</tr>
<tr>
<td>Blocks/cylinder</td>
<td>8029</td>
</tr>
<tr>
<td>Sectors/track</td>
<td>63</td>
</tr>
<tr>
<td>Heads</td>
<td>255</td>
</tr>
<tr>
<td>Spindle Speed</td>
<td>7200 rpm</td>
</tr>
<tr>
<td>Average Latency</td>
<td>4.16 msec</td>
</tr>
<tr>
<td>Random read seek time</td>
<td>&lt; 8.5 msec</td>
</tr>
<tr>
<td>Random read Write time</td>
<td>&lt; 9.5 msec</td>
</tr>
</tbody>
</table>
File Organizations

How do we organize records in a file?

• **Heap files**: records not in any particular order
  – Good for scans

• **Sorted files**: records sorted by particular fields
  – Scans in the sorted order or range scans in the sorted order

• **Indexes**: Data structures to organize records via trees or hashing.
  – Like sorted files, they speed up searches for a subset of records, based on values in certain (“search key”) fields
  – Updates are much faster than in sorted files
Comparing File Organizations

Consider an employee table with search key <age,sal>

- **Scans**: fetch all records in the file
- **Point queries**: find all employees who are 30 years old (let’s assume there’s only one such employee)
- **Range queries**: find all employees aged above 65.
- **Insert** a record.
- **Delete** a record given its RID.
Analysis of Algorithms

• Computation model
  – CPU comparison operation
  – General: most expensive operation

• Worst-case
  – How bad can it get?

• Average-case
  – Assumption about probabilities

• Analysis: count the number of some operation w.r.t. some input size

• Asymptotics: Big “O”
  – Constants don’t matter
  – $500n + 10000 = O(n)$

SELECT *
FROM Employees E
WHERE E.age=30

For each tuple t in Employees {
  if (t.age==30)
  {
    output t
  }
}

Assume input size: $n$ tuples

What is the worse case number of output tuples?

What is the worse case running time in the number of comparisons?
**Search Algorithms on Sorted Data**

**Shortcircuited Linear Search**

For each tuple \( t \) in Employees

\[
\begin{align*}
&\text{if } (t.\text{age} == 30) \\
&\quad \text{output } t \\
&\text{elsif } (t.\text{age} > 30) \\
&\quad \text{exit}
\end{align*}
\]

**Binary Search**

\[
(lo, hi) = (0, n-1)
\]

\[
\text{mid} = lo + (hi - lo) / 2
\]

While \((hi > lo \&\& E[mid].\text{age} != 30)\)

\[
\begin{align*}
&\text{if } (E[mid].\text{age} < 30) \\
&\quad \text{lo} = \text{mid} \\
&\text{else} \\
&\quad \text{hi} = \text{mid} \\
&\quad \text{mid} = lo + (hi - lo) / 2
\end{align*}
\]

Output all satisfying tuples around \( E[mid] \)

What is the worse case running time in the number of comparisons ?

SELECT *
FROM Employees E
WHERE E.age=30

Tuples are sorted by age
Analysis of Binary Search

- Number tuples searched per iteration = n, n/2, n/4, ... 1
- Hence the number of iterations = O( log n )
- Therefore number of comparisons = O( log n )

(lo, hi) = (0, n-1)
mid = lo + (hi - lo)/2
While(hi>lo && E[mid].age!=30)
{
  if (E[mid].age < 30)
    {lo=mid}
  else
    {hi=mid}
  mid = lo + (hi - lo)/2
}
Output all satisfying tuples around E[mid]
Analysis of DBMS Algorithms

```
SELECT *
FROM Employees
WHERE age=30
```

for each page \( p \) of Employees table
{
    if (\( p \) not in bufferpool)
    {
        Fetch \( p \) from disk
    }
    for each tuple \( t \) in page \( p \)
    {
        if (\( t \).age==40)
        {
            output \( t \)
        }
    }
}

Worst case running time =
+ time to fetch all pages of Employees from disk
+ time to compare age
+ time to output result

Table Scan

How would you estimate these times?

What is the worst case number of disk access?

What is the most expensive operation?
Analysis Model

• B : number of data pages
• R : number of records per page
• D : average time to read/write a disk page
  – From previous calculations, if a page is 2K bytes, D is about 13 milliseconds
• C : average time to process a record
  – For the 1 Ghz processors we have today, assuming it takes 100 cycles, C is about 100 nanoseconds
Table Scans on Heap Files

for each page $p$ of Employees table
{
  if ($p$ not in bufferpool)
  {
    Fetch $p$ from disk
  }
  for each tuple $t$ in page $p$
  {
    output $t$
    if ($t$.age==30)
    {
      output $t$
    }
    if ($t$.age>20 && $t$.age<30)
    {
      output $t$
    }
  }
}

SELECT *
FROM Employees
WHERE age=30

O(B) pages get fetched +
O(B*R) tuples processed

SELECT *
FROM Employees
WHERE age > 20 and age < 30

SELECT *
FROM Employees