ICS 321 Fall 2012
Algebraic and Logical Query Languages (ii)

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• A (relational) **atom**
  – Consists of a **predicate** and a list of **arguments**
  – Arguments can be **constants** or **variables**
  – Takes on **Boolean value** (true or false)

• A relation R can be represented as a predicate R
  – A tuple \(<a, b, c, d, e, f, g>\) is in R iff the atom R(a, b, c, d, e, f, g) is true.
Example: tables in datalog

\[
\begin{array}{|c|c|}
\hline
A & B \\
\hline
1 & 2 \\
3 & 4 \\
\hline
\end{array}
\]

R(1,2) and R(3,4) are true by default. R(1,4) would be false.
Arithmetic Atoms

\[ x < y \]
\[ x + 1 \geq y + 4z \]

Can contain both constants and variables.
Datalog Rules

LongMovie(t,y) :- Movies(t,y,l,g,s,p), l >= 100

(t,y) is a tuple of LongMovie
IF (t,y,l,g,s,p) is a tuple of Movies and length of movie is at least 100

These two “t,y” have to match
These two “l” have to match

Aka “subgoal”
Can be preceded by negation operator “NOT” or “~”

Anonymous variables
Every **variable** that appears anywhere in the rule **must** appear in some **nonnegated, relational subgoal** of the body

- Without the safety condition, rules may be underspecified, resulting in an infinite relation (not allowed).

- Examples
  - LongMovie(t,y) :- Movies(t,y,l,_,_,_,) , l >=100
  - P(x,y) :- Q(x,z), NOT R(w,x,z), x<y
Alternative Interpretation: Consistency

For each consistent assignment of nonnegated, relational subgoal,

Check the negated, relational subgoals and the arithmetic subgoals for consistency

<table>
<thead>
<tr>
<th>Q(x,z)</th>
<th>R(z,y)</th>
<th>Consistent?</th>
<th>NOT Q(x,y)</th>
<th>Head</th>
</tr>
</thead>
<tbody>
<tr>
<td>(1,2)</td>
<td>(2,3)</td>
<td>Yes</td>
<td>false</td>
<td></td>
</tr>
<tr>
<td>(1,2)</td>
<td>(3,1)</td>
<td>No, z=2,3</td>
<td></td>
<td></td>
</tr>
<tr>
<td>(1,3)</td>
<td>(2,3)</td>
<td>No, z=2,3</td>
<td></td>
<td></td>
</tr>
<tr>
<td>(1,3)</td>
<td>(3,1)</td>
<td>Yes</td>
<td>true</td>
<td>P(1,1)</td>
</tr>
</tbody>
</table>
Intensional vs Extensional

- **Extensional** predicates – relations stored in a database
- **Intensional** predicates – computed by applying one or more datalog rules

**Datalog**

\[
\begin{align*}
Q(1,2) \\
Q(1,3) \\
R(2,3) \\
R(3,1) \\
P(x,y) & : Q(x,z), R(z,y), \text{NOT } Q(x,y)
\end{align*}
\]
What about bag semantics?

• Datalog still works if there are no negated, relational subgoals.
• Treat duplicates like non-duplicates

Datalog

R(1,2)  R(1,2)  S(2,3)  S(4,5)  S(4,5)  H(x,z) :- R(x,y), S(y,z)

<table>
<thead>
<tr>
<th>R(x,y)</th>
<th>S(y,z)</th>
<th>Consistent?</th>
<th>Head</th>
</tr>
</thead>
<tbody>
<tr>
<td>(1,2)</td>
<td>(2,3)</td>
<td>Yes</td>
<td>H(1,3)</td>
</tr>
<tr>
<td>(1,2)</td>
<td>(4,5)</td>
<td>No, y=2,4</td>
<td></td>
</tr>
<tr>
<td>(1,2)</td>
<td>(4,5)</td>
<td>No, y=2,4</td>
<td></td>
</tr>
<tr>
<td>...</td>
<td>...</td>
<td>...</td>
<td>...</td>
</tr>
</tbody>
</table>
Example 1

**Datalog**

Answer(x,y) :- A(x,y)
Answer(x,y) :- B(x,y)
Example 2

Datalog

Answer(x,y) :- A(x,y), B(x,y)
Example 3

Datalog

Answer(x,y) :- A(x,y), NOT B(x,y)
Example 4

**Datalog**

\[
\text{Answer}(x,y) \ :- \ A(x,y), \ x > 10, \ y = 200
\]
Example 5

**Datalog**

Answer(x) :- A(x,y)
Example 6

\[
\text{Datalog} \\
\text{Answer}(w,x,y,z) :\ A(w,x), B(y,z)
\]
Example 7

Datalog

Answer(w,x,y) :- A(w), B(x,y)
Example 8

**Datalog**

Answer(w,x,z) :- A(w,x), B(y,z), x>y
Example 9

\textbf{Datalog}

\begin{align*}
\text{Path}(x,y) & :\text{-} \text{Edge}(x,y) \\
\text{Path}(x,z) & :\text{-} \text{Edge}(x,y), \text{Edge}(y,z)
\end{align*}